

# A(nother) Random Variable Monad

## Preliminary Report

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## A Disclaimer

- ▶ This is *NOT* a talk about stochastic processes or Skorohod's Theorem.

Apologies – a funny thing happened on the way to the talk...

- ▶ Instead, I'll talk about some preliminary results on computational models for probabilistic choice.

Active theme over last several years

Other talks at the meeting: Plotkin, Barker, Jung

## Some History

- ▶ The *probabilistic powerdomain* is the traditional domain model for probabilistic choice.
  - ▶ Realized on DCPO as  $\mathbb{V}(D)$  – the space of Scott-continuous valuations of  $\mathcal{O}(D)$
  - ▶ Defines a monad on DCPO
  - ▶ Leaves Coh invariant, but  
*no CCC of domains is known to be invariant*
  - ▶ No distributive law wrt power domains
- ▶ Varacca (2003) devised a monad that overcomes these problems
  - ▶ Weakens the probabilistic law:  $p +_r p \neq p$ .
- ▶ Based on Varacca's work, M. (2007) devised a monad of *finite random variables* that leaves RB and FS invariant, and satisfies a distributive law.

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- ▶ Based on Varacca's work, M. (2007) devised a monad of *finite random variables* that leaves RB and FS invariant, and satisfies a distributive law.
- ▶ Goubault-Larrecq and Varacca (2011) proposed a monad of *continuous random variables*, but the Kleisli lift failed to be Scott continuous.
- ▶ Their work inspired this work – and that by Tyler Barker you'll hear after this talk.

## Some Background

- ▶ *Random variables versus Prob*

- ▶ Order on  $Prob(D)$  is more complicated than that on  $D$ :

$$\mu \sqsubseteq \nu \quad \text{iff} \quad \mu(U) \leq \nu(U) \quad (\forall U \text{ open})$$

$$d \mapsto \delta_d: D \rightarrow Prob(D) \text{ unit of monad}$$

- ▶ *Idea of Random Variables:*

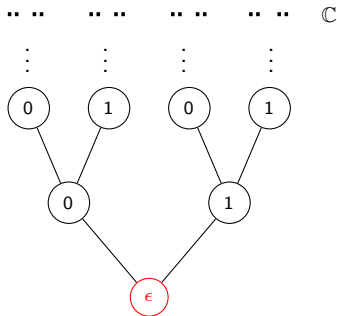
Choose fixed domain  $\mathcal{C}$  for which  $Prob(\mathcal{C})$  is well-behaved,  
and use

$f: \mathcal{C} \rightarrow D$  to define outcomes of probabilistic choices in  $\mathcal{C}$

## Some Background

- ▶ *Random variables versus Prob*
  - ▶ *Which domain  $\mathcal{C}$  to choose?*

Use the Cantor tree – the full binary tree together with the Cantor set at the top:



- ▶ So, the model looks like  $Prob(\mathcal{C}) \times [\mathcal{C} \rightarrow D]$  for a domain  $D$

## Sanity Check

$RanV(D) \equiv Prob(C) \times [C \rightarrow D]$  defines a functor:

$$g: D \rightarrow E \implies RanV(g)(\mu, f) = \langle \mu, g \circ f \rangle.$$

## Sanity Check

But to be a monad, we need a Koisli lift:

$$\begin{array}{ccc} \text{Prob}(\mathcal{C}) \times [\mathcal{C} \rightarrow D] & & \\ \eta_D \uparrow & \searrow h^\dagger & \\ D & \xrightarrow{h} & \text{Prob}(\mathcal{C}) \times [\mathcal{C} \rightarrow E] \end{array}$$

where  $\eta_D(d) = \langle \delta_\epsilon, \text{const}_d \rangle$ .

If  $h^\dagger(\mu, f) = \langle \nu, g \rangle$ ,

$g$  is easy:  $f: \mathcal{C} \rightarrow D$  &  $\pi_2 \circ h: D \rightarrow E \Rightarrow g = \pi_2 \circ h \circ f: \mathcal{C} \rightarrow E$ .



## Sanity Check

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where  $\eta_D(d) = \langle \delta_\epsilon, \text{const}_d \rangle$ .

If  $h^\dagger(\mu, f) = \langle \nu, g \rangle$ ,

For  $\nu$ , we should combine  $\mu$  and the family  $\{\pi_1 \circ h \circ f(d)\}_{d \in \text{supp } \mu}$ .

Suppose  $\mu = \sum_{d \in F} r_d \delta_d$  is simple. Then we expect:

$$h^\dagger \left( \sum_{d \in F} r_d \delta_d, f \right) = \left\langle \sum_{d \in F} r_d (\delta_d \otimes \pi_1 \circ h \circ f(d)), \pi_2 \circ h \circ f \right\rangle$$

where  $\otimes$  is a composition operation on measures.

$\otimes$  also should also be monotone!

## More Background

- ▶ Prob forms a monad on:
  - ▶ Comp – category of compact  $T_2$  spaces and continuous maps  
Algebras are compact affine spaces
  - ▶ CompMon – category of compact monoids and continuous monoid morphisms  
Algebras are compact affine monoids
  - ▶ CompGrp – category of compact groups and continuous group maps  
Algebras again compact affine monoids – right adjoint changes to  $M \mapsto H(1_M)$
- ▶ On CompMon, the monoid operation on  $Prob(S, \cdot, e)$  is *convolution*:  
 $*$ :  $Prob(S) \times Prob(S) \rightarrow Prob(S)$  by

$$\mu * \nu(A) = \cdot (\mu \times \nu)(A) = \mu \times \nu(\cdot^{-1}(A)) = \mu \times \nu(\{\langle x, y \rangle \mid x \cdot y \in A\}).$$

## More Background

- ▶ An *ordered monoid* is a monoid  $(S, \cdot, e)$  which has a partial order relative to which  $\cdot : S \times S \rightarrow S$  is monotone.
  - ▶  $\mathcal{C}$  is a monoid under concatenation, but it is *not* an ordered monoid.
  - ▶  $\mathcal{C}_\surd \equiv \{0, 1\}^\infty \cup \{0, 1\}^* \surd$  with the prefix order, and let
$$s \cdot t = \begin{cases} s't & \text{if } s = s'\surd \\ s & \text{otherwise} \end{cases}$$
  - ▶  $(\mathcal{C}_\surd, \cdot, \surd)$  is a compact ordered monoid.
- ▶ Prob is a monad on  $\text{CompOrdMon}$  – category of compact, ordered monoid and continuous, monotone monoid morphisms.

The convolution operation on  $\text{Prob}(S, \cdot, e, \leq)$  is monotone:

$$\mu \sqsubseteq \mu', \nu \sqsubseteq \nu' \implies \mu * \nu \sqsubseteq \mu' * \nu'.$$

- ▶ This applies to *any* bounded complete domain in the Lawson topology – in particular to  $\mathcal{C}_\surd$ .

## Towards a Monad

The monad should be:  $RanV(D) = Prob(\mathcal{C}_{\sqrt{}}) \times [\mathcal{C}_{\sqrt{}} \rightarrow D]$  where

$$\begin{array}{ccc} Prob(\mathcal{C}_{\sqrt{}}) \times [\mathcal{C}_{\sqrt{}} \rightarrow D] & & \\ \eta_D \uparrow & \searrow h^\dagger & \\ D & \xrightarrow{h} & Prob(\mathcal{C}_{\sqrt{}}) \times [\mathcal{C}_{\sqrt{}} \rightarrow E] \end{array}$$

where:

$$\eta_D = \langle \delta_{\sqrt{}}, const_d \rangle \text{ and for } \mu = \sum_{d \in F} r_d \delta_d$$

$$h^\dagger(\mu, f) = \left\langle \sum_{d \in F} r_d (\delta_d * (\pi_1 \circ h \circ f(d))), \pi_2 \circ h \circ f \right\rangle$$

## Towards a Monad

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ToDo's:

- ▶ Validate monad structure – May require restricting  $\mu$  to be “thin”
- ▶ Work out equational laws

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ToDo's:

- ▶ Devise closed form for probability component with  $\mu$  non-discrete

$$\mu = \lim_n \sum_{x \in F_n} r_x \delta_x \implies \mu " * " \nu = \lim_n \sum_{x \in F_n} r_x (\delta_x * (\pi_1 \circ h \circ f(x)))$$