Probabilistic Programming and a Domain-theoretic Approach to Skorohod's Theorem

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Outline

- I. Stochastic Processes and Skorohod's Theorem
- **II. Domains and Probability Measures**
- III. Proving Skorohod's Theorem

Along the way, compare with the computational Polish space approach

Stochastic Processes and Skorohod's Theorem

A stochastic process is a time-indexed family $\{X_t \mid t \in T \subseteq \mathbb{R}_+\}$ of random variables / elements $X_t \colon \Omega \to S$, where $(\Omega, \Sigma_{\Omega}, \mu)$ is a probability space, and S is a Polish space.

Fact: If S is Polish, then so is $(Prob(S), d_p)$, where d_p is the *Prokhorov* metric.

Skorohod's Theorem

Let S be a Polish space, let $\nu \in \operatorname{Prob} S$, and let λ denote Lebesgue measure on [0,1]. Then there is a random variable $X : [0,1] \to S$ satisfying $X_* \lambda = \nu$.

Moreover, if $\nu_n, \nu \in \operatorname{Prob} S$ satisfy $\nu_n \to_w \nu$, then the random variables $X_n, X \colon [0,1] \to S$ with $X_* \lambda = \nu, X_{n*} \lambda = \nu_n$ satisfy $X_n \to X \lambda$ -a.e.

Fact: We could use any *standard probability space* (S, Σ_S, m) instead of $([0, 1], \lambda)$.

Domains

Domains are partially ordered sets with additional properties;

Informatic partial order

 $p \sqsubseteq q$ if q contains more information than p.

Example: Upper Space $U(X) = (\{K \subseteq X \mid \emptyset \neq K \text{ compact}\} \cup \{X\}, \supseteq).$

Directed completeness

 $\emptyset \neq D \subseteq P$ directed if $x, y \in D \Rightarrow (\exists z \in D) x, y \leq z$. *P* directed complete: *D* directed \Rightarrow sup *D* exists.

 $\mathcal{F} \subseteq U(X)$ directed $\Rightarrow \sup \mathcal{F} = \bigcap \mathcal{F}$.

Approximation

 $\begin{aligned} x \ll y \text{ iff } y \leq \sup D \implies (\exists d \in D) x \leq d. \\ Domain: \downarrow y = \{x \mid x \ll y\} \text{ directed and } y = \sup \downarrow y \\ K \ll L \text{ iff } L \subseteq K^{\circ}; \qquad L = \bigcap \{K \mid L \subseteq K^{\circ}\} = \sup \{K \mid K \ll L\}. \end{aligned}$

Domains

Scott Topology

U Scott open if:

- $U = \uparrow U = \{x \in P \mid (\exists u \in U) \ u \le x\}$ and
- D directed, sup $D \in U \Rightarrow D \cap U \neq \emptyset$.

Morphisms

- $f: P \rightarrow Q$ is *Scott continuous* if:
- f is monotone, and
- $D \text{ directed} \Rightarrow f(\sup D) = \sup f(D).$

 $f: X \to Y \Rightarrow U(f): U(X) \to U(Y)$ by U(f)(K) = f(K)

is monotone and $f(\bigcap \mathcal{F}) = \bigcap f(\mathcal{F})$.

Lawson Topology

Basis: { $\uparrow x \setminus \uparrow F \mid F \in \mathcal{P}_{<\omega}D$ }

Hausdorff refinement of Scott topology.

All the domains we discuss are Lawson compact.

Domain Environments

Embedding X in $U(\overline{X})$

 $X \text{ Polish} \ \Rightarrow \ X \hookrightarrow [0,1]^{\omega} \ \Rightarrow \ X \subseteq_{G_{\delta}} \overline{X} \text{ compact Polish}$

 $\{K_n \mid K_n \in U(\overline{X}), n > 0\}$: neighborhood basis of compact subsets of \overline{X} . Then:

1°
$$X \subseteq \overline{X} \hookrightarrow \operatorname{Max} U(\overline{X}) \subseteq U(\overline{X})$$
 by $x \mapsto \{x\} = \bigcap_n \{K_n \mid x \in K_n^\circ\}$

- 2° X inherits the Scott topology = Lawson topology on Max $U(\overline{X})$.
- 3° Each family $\{K_1, \ldots, K_{n_m}\}$ with $\overline{X} \subseteq \bigcup_{n_m} K_i^{\circ}$ defines a

Scott-continuous projection $\psi_m \colon U(\overline{X}) \to U(\overline{X})$ with finite image.

4° Ordering the covers by refinement yields $\mathbf{1}_{U(\overline{X})} = \sup_{m} \psi_{m}$.

Domains and Probability Measures

Prob(D) is a Domain

- D (Lawson compact) domain \Rightarrow Prob(D) (Lawson compact) domain:
- 1° $\mu \leq \nu$ iff $\int f d\mu \leq \int f d\nu$ ($\forall f : D \to \mathbb{R}_+$ Scott continuous)
- 2° D Lawson compact \Rightarrow (Prob(D), weak) = (Prob(D), Lawson).

3°
$$D = U(\overline{X}), X$$
 Polish $\Rightarrow \mathbf{1}_{\operatorname{Prob}(D)} = \sup_{n} \psi_{n*}, \text{ so}$
 $\mu = \sup_{n} \psi_{n*} \mu, \text{ with } \psi_{n*} \mu = \sum_{i \leq m_n} \mu(K_i) \delta_{K_i} (\forall \mu).$

- 4° Choosing $r_i < \mu(K_i)$ & $L_i \ll K_i$ and $\epsilon = 1 (\sum_i r_i)$ produces $\nu_n \stackrel{\text{def}}{=} \epsilon \delta_X + \sum_{i \le m_n} r_i \delta_{L_i} \ll \psi_{n*} \mu.$
- 5° Construct $\nu_1 \ll \cdots \ll \nu_n \ll \nu_{n+1} \ll \cdots$ with $\mu = \sup_n \nu_n$.

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By comparison:

 $\begin{array}{ll} (X,S,d) \text{ a computational Polish space implies} \\ \mu = \lim_{n} \sum_{i \leq m_n} d_i \delta_{x_i} \in \operatorname{Max} U(\overline{X}) \text{ with } x_i \in S. \\ \end{array}$ $\begin{array}{ll} \text{Then} & \nu_n \ll \psi_{n*} \mu \leq \mu \ \Rightarrow (\exists N) \ n \geq N \ \Rightarrow \ \nu_n \ll \sum_{i \leq m_n} d_i \delta_{x_i}. \end{array}$

A Domain Environment for $\ensuremath{\mathcal{C}}$



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• Given
$$\mu \in \operatorname{Prob} U(\overline{X})$$
 and $\psi_{n*}\mu = \sum_{i \leq m_n} \mu(K_i)\delta_{K_i}$, fix p_n and choose $r_i \in D_{p_n} = \{\frac{s}{2^{p_n}} \mid 0 \leq s \leq 2^{p_n}\}$ and $L_i \ll K_i$.

• Then
$$\nu_n \stackrel{\text{def}}{=} (1 - \sum_i r_i) \delta_X + \sum_{i \le m_n} r_i \delta_{L_i} \ll \sum_{i \le m_n} \mu(\kappa_i) \delta_{\kappa_i}$$
.

• Define
$$f_n : \mathcal{C}_{p_n} \to U(\overline{X})$$
 by $f_n(j) = \begin{cases} L_1 & \text{if } 1 \le j \le r_1 \\ L_2 & \text{if } r_1 < j \le r_1 + r_2 \\ \vdots \\ X & \text{if } \sum_i r_i < j \end{cases}$

Then $f_{n*}(\frac{1}{2^{p_n}}\sum_{i\leq 2^{p_n}}\delta_{\frac{i}{2^{p_n}}})=(1-\sum_i r_i)\delta_X+\sum_{i\leq m_n}r_i\delta_{L_i}=\nu_n.$

• So, $f_n \circ \pi_n \colon \mathcal{C} \to U(\overline{X})$ is Lawson continuous and $(f_n \circ \pi_n)_* (\mu_{\mathcal{C}}) = \nu_n = (1 - \sum_i r_i)\delta_X + \sum_{i \leq m_n} r_i \delta_{L_i}.$

A Domain Environment for C

- Given $\mu \in \operatorname{Prob} U(\overline{X})$ and $\psi_{n*}\mu = \sum_{i \leq m_n} \mu(K_i)\delta_{K_i}$, fix p_n and choose $r_i \in D_{p_n} = \{\frac{s}{2^{p_n}} \mid 0 \leq s \leq 2^{p_n}\}$ and $L_i \ll K_i$.
- Then $\nu_n \stackrel{\text{def}}{=} (1 \sum_i r_i) \delta_X + \sum_{i \le m_n} r_i \delta_{L_i} \ll \sum_{i \le m_n} \mu(K_i) \delta_{K_i}$.

Then
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- So, $f_n \circ \pi_n \colon \mathcal{C} \to U(\overline{X})$ is Lawson continuous and
- $(f_n \circ \pi_n)_* (\mu_{\mathcal{C}}) = \nu_n = (1 \sum_i r_i)\delta_X + \sum_{i \leq m_n} r_i \delta_{L_i}.$
- By construction $f_n : \downarrow C_{P_n} \to U(\overline{X})$ satisfies $f_n \leq f_{n+1}$, and then $X = (\sup_n f_n \circ \pi_n)|_{\mathcal{C}} : \mathcal{C} \to U(\overline{X})$ measurable with $X_* \mu_{\mathcal{C}} = \mu$.

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• Argue directly that $X_m \to X$ a.s. λ .

Questions?